

Solving Tropical Optimization Problems via Matrix Sparsification

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Introduction: Tropical Optimization

- ▶ **Tropical (idempotent) mathematics** focuses on the theory and applications of semirings with idempotent addition
- ▶ **The tropical optimization problems** are those that are formulated and solved within the framework of tropical mathematics
- ▶ Many problems have **objective functions** defined on vectors over **idempotent semifields** (semirings with multiplicative inverses)
- ▶ Both unconstrained problems and problems with constraints in the form of **linear inequalities and equalities** are considered
- ▶ The problems find **applications** in various areas, including
 - ▶ project scheduling,
 - ▶ location analysis,
 - ▶ transportation networks,
 - ▶ decision making,
 - ▶ discrete event systems

Idempotent Algebra: Definitions and Notation

Idempotent Semifield

- ▶ *Definition:* the algebraic system $\langle \mathbb{X}, 0, \mathbb{1}, \oplus, \otimes \rangle$
- ▶ *Carrier set:* \mathbb{X} with neutral elements, *zero* 0 and *identity* $\mathbb{1}$
- ▶ *Associative and commutative* binary operations: \oplus and \otimes
- ▶ Multiplication \otimes is *distributive* over addition
- ▶ Addition \oplus is *idempotent*: $x \oplus x = x$ for all $x \in \mathbb{X}$
- ▶ Multiplication is *invertible*: for each nonzero $x \in \mathbb{X}$, there exists an inverse $x^{-1} \in \mathbb{X}$ such that $x \otimes x^{-1} = \mathbb{1}$
- ▶ *Linear order:* the order $x \leq y \iff x \oplus y = y$ is a total order
- ▶ *Algebraic completeness:* the equation $x^p = a$ is solvable for any $a \in \mathbb{X}$ and integer p to provide powers with rational exponents
- ▶ *Notational convention:* the multiplication signs \otimes will be omitted

Semifield $\mathbb{R}_{\max,+}$ (Max-Plus Algebra)

- ▶ **Definition:** $\mathbb{R}_{\max,+} = \langle \mathbb{R} \cup \{-\infty\}, -\infty, 0, \max, + \rangle$
- ▶ **Carrier set:** $\mathbb{X} = \mathbb{R} \cup \{-\infty\}$; **zero and identity:** $0 = -\infty$, $1 = 0$
- ▶ **Binary operations:** $\oplus = \max$ and $\otimes = +$
- ▶ **Idempotent addition:** $x \oplus x = x$ for all x ($\max(x, x) = x$)
- ▶ **Multiplicative inverse:** for each $x \in \mathbb{R}$, there exists x^{-1} ($= -x$)
- ▶ **Power notation:** for each $x, y \in \mathbb{R}$, there is defined x^y ($= xy$)
- ▶ Further examples of real idempotent semifields:

$$\mathbb{R}_{\min,+} = \langle \mathbb{R} \cup \{+\infty\}, +\infty, 0, \min, + \rangle,$$

$$\mathbb{R}_{\max,\times} = \langle \mathbb{R}_+ \cup \{0\}, 0, 1, \max, \times \rangle,$$

$$\mathbb{R}_{\min,\times} = \langle \mathbb{R}_+ \cup \{+\infty\}, +\infty, 1, \min, \times \rangle,$$

where $\mathbb{R}_+ = \{x \in \mathbb{R} | x > 0\}$

Vector and Matrix Algebra Over \mathbb{X}

- ▶ The idempotent semifield \mathbb{X} is routinely extended to the idempotent systems of vectors in \mathbb{X}^n and of matrices in $\mathbb{X}^{m \times n}$
- ▶ The matrix and vector operations follow the usual entry-wise formulae with \oplus as addition, and \otimes as multiplication
- ▶ For any vectors $\mathbf{a} = (a_i)$ and $\mathbf{b} = (b_i)$ in \mathbb{X}^n , and a scalar $x \in \mathbb{X}$, the vector operations follow the conventional rules

$$\{\mathbf{a} \oplus \mathbf{b}\}_i = a_i \oplus b_i, \quad \{x\mathbf{a}\}_i = xa_i$$

- ▶ For any matrices $\mathbf{A} = (a_{ij}) \in \mathbb{X}^{m \times n}$, $\mathbf{B} = (b_{ij}) \in \mathbb{X}^{m \times n}$ and $\mathbf{C} = (c_{ij}) \in \mathbb{X}^{n \times l}$, and $x \in \mathbb{X}$, the matrix operations are given by

$$\{\mathbf{A} \oplus \mathbf{B}\}_{ij} = a_{ij} \oplus b_{ij}, \quad \{\mathbf{AC}\}_{ij} = \bigoplus_{k=1}^n a_{ik}c_{kj}, \quad \{x\mathbf{A}\}_{ij} = xa_{ij}$$

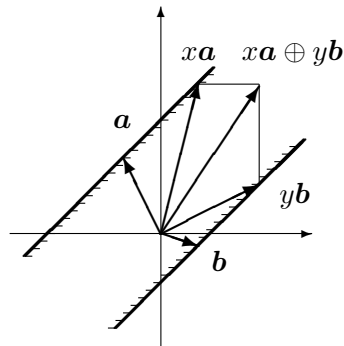
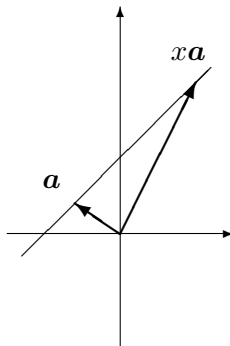
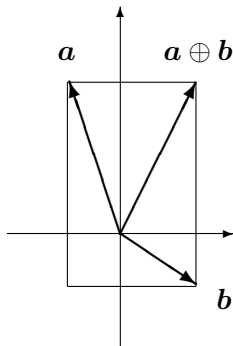
Idempotent Semimodule Over \mathbb{X}

- ▶ *Definition:* the system $\langle \mathbb{X}^n, \mathbf{0}, \oplus \rangle$ with scalar multiplication \otimes
- ▶ *Carrier set:* the set of column vectors of order n denoted \mathbb{X}^n
- ▶ *Zero vector:* $\mathbf{0} = (0, \dots, 0)^T$, *vector of ones:* $\mathbf{1} = (1, \dots, 1)^T$
- ▶ *Operations:* vector addition \oplus , and scalar multiplication \otimes
- ▶ *Regular vector:* any vector without zero components
- ▶ *Multiplicative conjugate transposition* transforms any nonzero column vector $x = (x_i)$ into the row vector $x^- = (x_i^-)$, where

$$x_i^- = \begin{cases} x_i^{-1}, & \text{if } x_i \neq 0; \\ 0, & \text{otherwise} \end{cases}$$

- ▶ *Linear dependence:* a vector y is linearly dependent on vectors x_1, \dots, x_m if $y = c_1 x_1 \oplus \dots \oplus c_m x_m$ for some scalars c_1, \dots, c_m

Graphical Representation for $\mathbb{R}_{\max,+}^2$



- ▶ Addition (left), scalar multiplication (middle), and a linear span (right) of vectors in the Cartesian coordinate system in the plane

Matrices Over \mathbb{X} (Further Definitions)

- ▶ *Zero* and *identity matrices*:

$$\mathbf{0} = \begin{pmatrix} 0 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0 \end{pmatrix}, \quad \mathbf{I} = \begin{pmatrix} 1 & & 0 \\ & \ddots & \\ 0 & & 1 \end{pmatrix}$$

- ▶ *Row- (column-) regular matrix*: any matrix without rows (columns) that consist entirely of zeros
- ▶ *Multiplicative conjugate transposition* transforms any nonzero matrix $\mathbf{A} = (a_{ij})$ into the matrix $\mathbf{A}^- = (a_{ij}^-)$, where

$$a_{ij}^- = \begin{cases} a_{ji}^{-1}, & \text{if } a_{ji} \neq 0; \\ 0, & \text{otherwise} \end{cases}$$

Tropical Optimization Problems: Examples

Linear Objective Functions

- ▶ Hoffman (1963), Superville (1978), U. Zimmermann (1981)

$$\begin{array}{ll} \text{minimize} & p^T x, \\ \text{subject to} & Ax \geq d \end{array} \quad (\text{direct solution})$$

- ▶ K. Zimmermann (1984, 1992, 2003, 2006)

$$\begin{array}{ll} \text{minimize} & p^T x, \\ \text{subject to} & Ax \leq d, \quad Cx \geq b, \\ & g \leq x \leq h \end{array} \quad (\text{algorithmic solution})$$

- ▶ Butkovič (1984, 2010), Butkovič and Aminu (2009)

$$\begin{array}{ll} \text{minimize} & p^T x, \\ \text{subject to} & Ax \oplus b = Cx \oplus d \end{array} \quad (\text{algorithmic solution})$$

Nonlinear Objective Functions

- ▶ Cuninghame-Green (1962, 1979), Engel and Schneider (1975), Elsner and van den Driessche (2004, 2010), K. (2013,2014)

$$\text{minimize } x^{-}Ax \quad (\text{direct solution})$$

- ▶ Cuninghame-Green (1976), U. Zimmermann (1981)

$$\begin{aligned} &\text{minimize } (Ax)^{-}d, \\ &\text{subject to } Ax \leq d \end{aligned} \quad (\text{direct solution})$$

- ▶ K. Zimmermann (1984)

$$\begin{aligned} &\text{minimize } (Ax)^{-}p \oplus p^{-}Ax, \\ &\text{subject to } g \leq x \leq h \end{aligned} \quad (\text{algorithmic solution})$$

Nonlinear Objective Functions

- ▶ K. (2004, 2009, 2012)

$$\text{minimize } (\mathbf{Ax})^{-} \mathbf{p} \oplus \mathbf{q}^{-} \mathbf{Ax} \quad (\text{direct solution})$$

- ▶ Butkovič and Tam (2009)

$$\text{minimize } \mathbf{1}^T \mathbf{Ax} (\mathbf{Ax})^{-} \mathbf{1}; \quad (\text{direct solution})$$

$$\text{maximize } \mathbf{1}^T \mathbf{Ax} (\mathbf{Ax})^{-} \mathbf{1} \quad (\text{direct solution})$$

- ▶ Gaubert, Katz and Sergeev (2012)

$$\begin{aligned} &\text{minimize } (\mathbf{p}^T \mathbf{x} \oplus r)(\mathbf{q}^T \mathbf{x} \oplus s)^{-1}, \\ &\text{subject to } \mathbf{Ax} \oplus \mathbf{b} \leq \mathbf{Cx} \oplus \mathbf{d} \end{aligned} \quad (\text{algorithmic solution})$$

Solution via Matrix Sparsification: Problem Formulation

Problem

Given a matrix $A \in \mathbb{X}^{m \times n}$ and vectors $p \in \mathbb{X}^m$, $q \in \mathbb{X}^n$, the problem is to find regular vectors $x \in \mathbb{X}^n$ that

$$\text{minimize } q^T x (Ax)^T p$$

- ▶ The problem appears in approximation in the sense of span seminorm (the maximum deviation between elements of a vector)
- ▶ Applications include project scheduling, decision making

Partial Solution

Problem

$$\text{minimize } q^- x (Ax)^- p$$

Proposition

Let A be a row-regular matrix, p be nonzero and q regular vectors.

Then, the minimum is $\Delta = (Aq)^- p$, and attained at $x = \alpha q$ for all $\alpha > 0$

Sketch of Proof.

1. $xx^- \geq I \implies (q^- x)^{-1} x = (q^- x x^-)^- \leq q$
2. $(q^- x)^{-1} x \leq q \implies (q^- x)^{-1} Ax \leq Aq$
3. $(q^- x)^{-1} Ax \leq Aq \implies q^- x (Ax)^- p \geq (Aq)^- p = \Delta$
4. $x = \alpha q \implies q^- x (Ax)^- p = (Aq)^- p = \Delta$ □

Problem

$$\text{minimize } q^- x (Ax)^- p$$

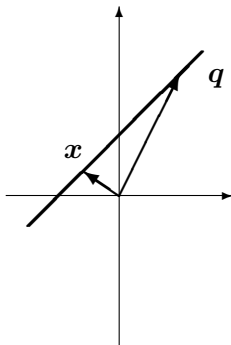
Example (in terms of $\mathbb{R}_{\max,+}$)

$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

Partial Solution

$$\Delta = (Aq)^- p = 2,$$

$$x = \alpha q, \quad \alpha \in \mathbb{R}$$



Characterization of Solution

Problem

$$\text{minimize } q^- x (Ax)^- p$$

Lemma

Let A be row-regular, p be nonzero, q be regular, and $\Delta = (Aq)^- p$.

Then, all regular solutions are given by

$$q^- x = \alpha, \quad Ax \geq \alpha \Delta^{-1} p, \quad \alpha > 0$$

Sketch of Proof.

1. $q^- x (Ax)^- p = \Delta \iff q^- x = \alpha, (Ax)^- p = \alpha^{-1} \Delta, \alpha > 0$
2. $(Ax)^- p = \alpha^{-1} \Delta \iff (Ax)^- p \leq \alpha^{-1} \Delta, (Ax)^- p \geq \alpha^{-1} \Delta$
3. $q^- x = \alpha \implies x \leq \alpha q \implies (Ax)^- p \geq \alpha^{-1} \Delta$
4. $(Ax)^- p \leq \alpha^{-1} \Delta \iff p \leq \alpha^{-1} \Delta Ax \iff Ax \geq \alpha \Delta^{-1} p$ □

Properties of Solution

Problem

$$\text{minimize } q^- x (Ax)^- p$$

Corollary

Let A be row-regular, p be nonzero, q be regular, and $\Delta = (Aq)^- p$.

Then, the set of regular solutions is closed under addition and scalar multiplication

Sketch of Proof.

Addition:

$$\begin{array}{l} q^- x = \alpha, \quad Ax \geq \alpha \Delta^{-1} p, \\ q^- y = \beta, \quad Ay \geq \beta \Delta^{-1} p \end{array} \implies \begin{array}{l} q^- (x \oplus y) = \alpha \oplus \beta, \\ A(x \oplus y) \geq (\alpha \oplus \beta) \Delta^{-1} p \end{array}$$

Scalar multiplication: analogously



Matrix Sparsification

Problem

$$\text{minimize } \mathbf{q}^{-} \mathbf{x} (\mathbf{A} \mathbf{x})^{-} \mathbf{p}$$

Lemma

Let $\mathbf{A} = (a_{ij})$ be a row-regular matrix, $\mathbf{p} = (p_i)$ be a nonzero vector, $\mathbf{q} = (q_j)$ be a regular vector, and $\Delta = (\mathbf{A} \mathbf{q})^{-} \mathbf{p}$.

Then, replacing the matrix \mathbf{A} by the sparsified matrix $\hat{\mathbf{A}} = (\hat{a}_{ij})$, where

$$\hat{a}_{ij} = \begin{cases} a_{ij}, & \text{if } a_{ij} \geq \Delta^{-1} p_i q_j^{-1}; \\ 0, & \text{otherwise;} \end{cases}$$

does not change the solution

Sketch of Proof.

The representation $\mathbf{q}^{-} \mathbf{x} = \alpha$, $\mathbf{A} \mathbf{x} \geq \alpha \Delta^{-1} \mathbf{p}$ yields that there may be terms $a_{ij} x_j$, which do not affect the left-hand side of the inequality \square

Problem

minimize $q^- x (Ax)^- p$

Example (in terms of $\mathbb{R}_{\max,+}$)

$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

Threshold and Sparsified Matrices

$$\Delta = 2, \quad \Delta^{-1} p q^- = \begin{pmatrix} 2 & 1 \\ -1 & -2 \end{pmatrix}, \quad \hat{A} = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}$$

Extended Solution

Problem

$$\text{minimize } q^- x (Ax)^- p$$

Lemma

Let A be *sparsified* row-regular, p be nonzero, q be regular, and $\Delta = (Aq)^- p$.

Then, any vector given by the condition

$$x = (I \oplus \Delta^{-1} A^- p q^-) u, \quad u > 0$$

is a solution of the problem

Sketch of Proof.

- $x = (I \oplus \Delta^{-1} A^- p q^-) u \iff \alpha \Delta^{-1} A^- p \leq x \leq \alpha q$
- $\alpha \Delta^{-1} A^- p \leq x \leq \alpha q \implies q^- x = \alpha, Ax \geq \alpha \Delta^{-1} p$



Problem

minimize $q^-x(Ax)^-p$

Example (in terms of $\mathbb{R}_{\max,+}$)

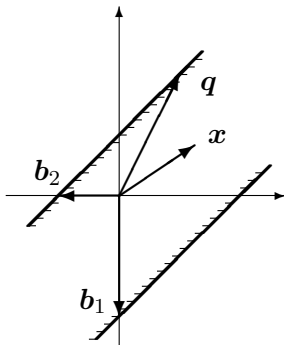
$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

Extended Solution

$$x = Bu, \quad u \in \mathbb{R}^2,$$

$$B = I \oplus \Delta^{-1} A^- p q^-$$

$$= \begin{pmatrix} 0 & -1 \\ -2 & 0 \end{pmatrix} = (b_1 \quad b_2)$$



All Solutions

Problem

$$\text{minimize } q^- x (Ax)^- p$$

Theorem

Let A be *sparsified* row-regular, p be nonzero, q be regular, and $\Delta = (Aq)^- p$.

Let \mathcal{A} be the set of matrices obtained from A by fixing one nonzero entry in each row and setting the others to $\mathbb{0}$.

Then, all regular solutions are given by

$$x = (I \oplus \Delta^{-1} A_1^- p q^-) u, \quad u > \mathbf{0}, \quad A_1 \in \mathcal{A}$$

Sketch of Proof.

$$x = (I \oplus \Delta^{-1} A_1^- p q^-) u, \quad A_1 \in \mathcal{A} \iff q^- x = \alpha, \quad Ax \geq \alpha \Delta^{-1} p \quad \square$$

Problem

minimize $q^- x (Ax)^- p$

Example (in terms of $\mathbb{R}_{\max,+}$)

$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

All Solutions

$$\Delta = 2, \quad A = \{A_1, A_2\}, \quad A_1 = \begin{pmatrix} 2 & 0 \\ 4 & 0 \end{pmatrix}, \quad A_2 = \begin{pmatrix} 2 & 0 \\ 0 & 1 \end{pmatrix};$$

$$x = B_1 u, \quad u \in \mathbb{R}^2, \quad B_1 = I \oplus \Delta^{-1} A_1^- p q^- = \begin{pmatrix} 0 & -1 \\ 0 & 0 \end{pmatrix};$$

$$x = B_2 u, \quad u \in \mathbb{R}^2, \quad B_2 = I \oplus \Delta^{-1} A_2^- p q^- = \begin{pmatrix} 0 & -1 \\ -2 & 0 \end{pmatrix}$$

Problem

minimize $q^- x (Ax)^- p$

Example (in terms of $\mathbb{R}_{\max,+}$)

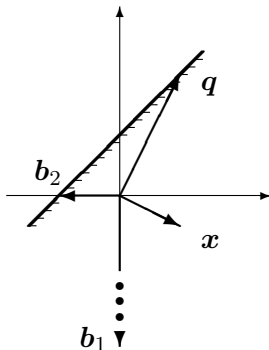
$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

All Solutions

$$x = B_1 u, \quad u \in \mathbb{R}^2,$$

$$B_1 = I \oplus \Delta^{-1} A_1^- p q^-$$

$$= \begin{pmatrix} 0 & -1 \\ 0 & 0 \end{pmatrix} = (b_1 \quad b_2)$$



Complete Solution

Problem

minimize $q^- x (Ax)^- p$

Theorem

Let A be *sparsified* row-regular, p be nonzero, q be regular, and $\Delta = (Aq)^- p$.

Let \mathcal{A} be the set of matrices obtained from A by leaving one nonzero entry in each row.

Let B_0 be the matrix whose columns form the maximal independent system of columns in matrices $B_1 = I \oplus \Delta^{-1} A_1^- p q^-$ for all $A_1 \in \mathcal{A}$.

Then, all regular solutions are given by $x = B_0 u$, $u > 0$

Sketch of Proof.

Follows from that the set of solutions to the problem is closed under vector addition and scalar multiplication □

Problem

minimize $q^- x (Ax)^- p$

Example (in terms of $\mathbb{R}_{\max,+}$)

$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

Complete Solution

$$B_1 = \begin{pmatrix} 0 & -1 \\ 0 & 0 \end{pmatrix}, \quad B_2 = \begin{pmatrix} 0 & -1 \\ -2 & 0 \end{pmatrix};$$

$$b_1 = \begin{pmatrix} 0 \\ 0 \end{pmatrix}, \quad b_2 = \begin{pmatrix} -1 \\ 0 \end{pmatrix}, \quad b_3 = \begin{pmatrix} 0 \\ -2 \end{pmatrix}, \quad b_3 = b_1 \oplus (-2)b_2;$$

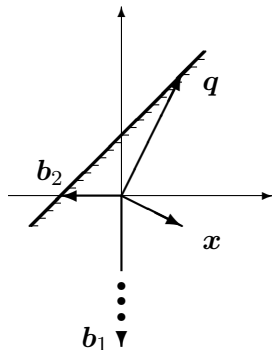
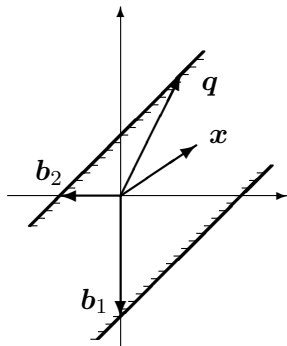
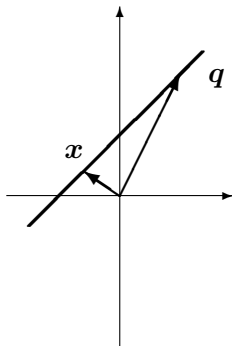
$$x = \begin{pmatrix} 0 & -1 \\ 0 & 0 \end{pmatrix} u, \quad u \in \mathbb{R}^2, \quad B_0 = B_1 = \begin{pmatrix} 0 & -1 \\ 0 & 0 \end{pmatrix}$$

Problem

minimize $q^- x (Ax)^- p$

Example (in terms of $\mathbb{R}_{\max,+}$)

$$A = \begin{pmatrix} 2 & 0 \\ 4 & 1 \end{pmatrix}, \quad p = \begin{pmatrix} 5 \\ 2 \end{pmatrix}, \quad q = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$



Concluding Remarks

- ▶ An optimization problem, which arises in approximation in the span seminorm, was examined in terms of tropical mathematics
- ▶ The problem is to minimize a function defined on vectors over an idempotent semifield by using conjugate transposition
- ▶ All solutions were characterized by simultaneous equation and inequality, and properties of the solution set were investigated
- ▶ A matrix sparsification technique was developed to derive a complete solution as a family of solution subsets
- ▶ The characteristic properties of solutions were exploited to describe the complete solution in a compact vector form
- ▶ The proposed solution approach can serve as a template to derive complete, direct solutions of other problems
- ▶ More solutions to tropical optimization problems with applications are available at http://arxiv.org/a/krivulin_n_1